

Not obviously manipulable allotment rules*

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Abstract

In the problem of allocating a single non-disposable commodity among agents whose preferences are single-peaked, we study a weakening of strategy-proofness called not obvious manipulability (NOM). If agents are cognitively limited, then NOM is sufficient to describe their strategic behavior. We characterize a large family of efficient and NOM rules, which we call "simple". The idea behind their definition has been dormant in the axiomatic literature for a long time. In economies with excess demand, simple rules fully satiate agents whose peak amount is no greater than equal division and assign, to each remaining agent, an amount between equal division and his peak. In economies with excess supply, simple rules are defined symmetrically. We also show that the single-plateaued domain is maximal for efficiency and NOM (together with other basic requirements). Therefore, even though replacing strategy-proofness with NOM greatly expands the family of admissible rules, the maximal domain of preferences involved remains basically unaltered.

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1 Introduction

Consider the problem of allocating a single non-disposable commodity among a group of agents with single-peaked preferences: up to some critical level, called the peak, an increase in an agent’s consumption raises his welfare; beyond that level, the opposite holds. In this context, an allotment rule is a systematic procedure that allows agents to select an allotment, among many, according to their preferences. An allotment rule is strategy-proof if misreporting preferences is never better than truth-telling. In this paper, we study a weakening of strategy-proofness called not obvious manipulability.

Strategy-proofness, together with efficiency and fairness, are considered the three perennial goals of mechanism design. The combination of these three goals singles out the “uniform” rule (Sprumont, 1991), whereas relaxing fairness by permitting an asymmetric treatment of agents expands the options to the versatile family of “sequential” allotment rules (Barberà et al., 1997). Insisting on efficiency, and in line with a current strand in the literature, we investigate what happens when strategy-proofness is weakened. The idea underlying this approach is that even though manipulations are pervasive, agents may not realize they can manipulate a rule because they lack information about others’ behavior or they are cognitively limited. A misreport that in a specific situation (a specific preference profile of the other agents) is a profitable manipulation may actually be worse than truth-telling in another situation. Therefore, if we consider an agent with limited information on other agents’ preferences, it may be very unclear whether such manipulation will be profitable in practice. Assuming that the agent knows all possible outcomes under any misreport and under truth-telling, there are several different ways to determine when a misreport is recognized or identified as a profitable manipulation. Troyan and Morrill (2020), in the context of two-sided matching, propose a simple and tractable way to do this by considering best/worst-case scenarios to formalize the notion of obvious manipulation. A manipulation is obvious if it either makes the agent better off than truth-telling in the worst case or makes the agent better off than truth-telling in the best case. An allotment rule is not obviously manipulable (NOM) if it has no obvious manipulation.

If agents are cognitively limited, then NOM is sufficient to describe their strategic behavior. Therefore, the question arises to what extent NOM rules enrich the landscape of strategy-proof rules. In order to keep our analysis tractable, we will consider own-peak-only rules. This means that the sole information collected by these rules from an agent’s

preference to determine his allotment is his peak amount. The own-peak-only property follows from efficiency and strategy-proofness (see [Sprumont, 1991](#); [Ching, 1994](#)). However, since we do not impose strategy-proofness, we explicitly invoke the property.

Our goal is to introduce a large family of own-peak-only and NOM rules, which we call "simple". Their definition is as follows. In economies with excess demand, simple rules fully satiate agents whose peak amount is no greater than equal division and assign to each remaining agent an amount between equal division and his peak. Symmetrically, in economies with excess supply, simple rules fully satiate agents whose peak is no smaller than equal division and assign to each remaining agent an amount between his peak and equal division. We believe that the full family of simple rules has been dormant in the literature because even though many of the rules (and families of rules) identified in the literature belong to this large family, their definition is more detailed and not as straightforward. This is the case, for example, of the family of sequential rules in [Barberà et al. \(1997\)](#) that are defined throughout an intricate process needed to cope with strategy-proofness.

Our main result is the characterization of the family of all simple rules. In [Theorem 2](#), we show that an allotment rule is simple if and only if it is own-peak-only, efficient, satisfies NOM, and meets the equal division guarantee. This last property is a minimal fairness requirement that states that whenever an agent demands equal division, the rule must guarantee him that amount. If we narrow the picture to symmetric rules, i.e., rules that assign indifferent allotments to agents with the same preferences, then we also obtain a characterization of a big subfamily of simple rules that satisfies all five properties ([Corollary 1](#)).¹ When we replace NOM with strategy-proofness, in contrast, the only efficient and symmetric rule that remains is the uniform rule ([Ching, 1994](#)). We also provide three variants of our main characterization: (i) changing own-peak-onliness by the more usual (and stronger) peaks-only requirement ([Theorem 3](#)), (ii) adapting the result to economies with individual endowments ([Theorem 4](#)), and (iii) invoking a peak monotonicity property that also encompasses symmetry and the equal division guarantee ([Theorem 5](#)).

Next, we analyze the maximality of the domain of preferences (including the domain of single-peaked preferences) for which own-peak-only, efficient, and NOM rules exist, meeting also the equal division guarantee. For the properties of efficiency, strategy-

¹Note that although symmetry and the equal division guarantee can both be considered as fairness requirements, they are logically independent properties.

proofness, and symmetry, the single-plateaued domain is maximal (Ching and Serizawa, 1998; Massó and Neme, 2001). In Theorem 6, we show that the single-plateaued domain is maximal for our properties as well. Therefore, even though replacing strategy-proofness with NOM greatly expands the family of admissible rules, the maximal domain of preferences involved remains basically unaltered.

To the best of our knowledge, our paper is the first one that applies Troyan and Morrill (2020) notion of obvious manipulations to the allocation of a non-disposable commodity among agents with single-peaked preferences. In the context of voting, Aziz and Lam (2021) and Arribillaga and Bonifacio (2022) study non-peaks-only and peaks-only rules, respectively. Other recent papers that study obvious manipulations, in other situations, are Ortega and Segal-Halevi (2022), Psomas and Verma (2022), and Arribillaga and Risma (2023).

The rest of the paper is organized as follows. The model and the concept of obvious manipulations are introduced in Section 2. In Section 3, we present the characterization of simple rules. Section 4 analyses some related characterizations. The maximal domain result is presented in Section 5. To conclude, some final remarks are gathered in Section 6.

2 Preliminaries

2.1 Model

A social endowment $\Omega \in \mathbb{R}_{++}$ is an amount of a perfectly divisible commodity to be distributed among a set of agents $N = \{1, 2, \dots, n\}$. Each $i \in N$ is equipped with a continuous preference relation R_i defined over $\mathbb{R}_+ \cup \{\infty\}$. Call P_i and I_i to the strict preference and indifference relations associated with R_i , respectively. Denote by \mathcal{U} the domain of all such preferences. Given $R_i \in \mathcal{U}$, let $p(R_i) = \{x \in \mathbb{R}_+ \cup \{\infty\} : x R_i y \text{ for each } y \in \mathbb{R}_+ \cup \{\infty\}\}$ be the set of preferred consumptions according to R_i , called the **peak** of R_i . When $p(R_i)$ is a singleton, we slightly abuse notation and use $p(R_i)$ to denote its single element.

A (generic) domain of preferences \mathcal{D} is a subset of \mathcal{U} . Given a domain of preferences $\mathcal{D} \subseteq \mathcal{U}$, an **economy** in \mathcal{D} consists of a profile of preferences $R = (R_j)_{j \in N} \in \mathcal{D}^N$ and a social endowment $\Omega \in \mathbb{R}_{++}$ and is denoted by (R, Ω) . Let $\mathcal{E}_{\mathcal{D}}$ be the domain of all such economies. Given a social endowment $\Omega \in \mathbb{R}_{++}$, the set of **(feasible) allocations** of Ω is $X(\Omega) = \{x \in \mathbb{R}_+^N : \sum_{j \in N} x_j = \Omega\}$. An **(allotment) rule** on $\mathcal{E}_{\mathcal{D}}$ is a function $\varphi : \mathcal{E}_{\mathcal{D}} \rightarrow \mathbb{R}_+^N$

such that $\varphi(R, \Omega) \in X(\Omega)$ for each $(R, \Omega) \in \mathcal{E}_{\mathcal{D}}$.

Given a rule φ defined on a generic domain $\mathcal{E}_{\mathcal{D}}$, some desirable properties we consider are listed next.

Efficiency: For each $(R, \Omega) \in \mathcal{E}_{\mathcal{D}}$, there is no $x \in X(\Omega)$ such that $x_i R_i \varphi_i(R, \Omega)$ for each $i \in N$ and $x_i P_i \varphi_i(R, \Omega)$ for some $i \in N$.

Efficiency is the usual Pareto optimality criterion. Under this condition, for each economy, the allocation selected by the rule should be such that there is no other allocation that all agents find at least as desirable and at least one agent (strictly) prefers.

A rule is strategy-proof if, for each agent, truth-telling is always optimal, regardless of the preferences declared by the other agents. To define it formally, let $R_i, R'_i \in \mathcal{D}$, and $\Omega \in \mathbb{R}_{++}$. Preference R'_i is a **manipulation of φ at (R_i, Ω)** if there is $R_{-i} \in \mathcal{D}^{n-1}$ such that $\varphi_i(R'_i, R_{-i}, \Omega) P_i \varphi_i(R_i, R_{-i}, \Omega)$.

Strategy-proofness: For each $i \in N$ and each $(R_i, \Omega) \in \mathcal{D} \times \mathbb{R}_{++}$, there is no manipulation of φ at (R_i, Ω) .

The following is an informational simplicity property stating that if an agent unilaterally changes his preference for another one with the same peak, then his allotment remains unchanged.²

Own-peak-onliness: For each $(R, \Omega) \in \mathcal{E}_{\mathcal{D}}$, each $i \in N$, and each $R'_i \in \mathcal{D}$ such that $p(R'_i) = p(R_i)$, we have $\varphi_i(R, \Omega) = \varphi_i(R'_i, R_{-i}, \Omega)$.

Analyzing the uniform rule, [Sprumont \(1991\)](#) derives the own-peak-only property from efficiency and strategy-proofness (see also [Ching, 1994](#)). Since we do not impose strategy-proofness, in order to keep our analysis tractable we explicitly invoke own-peak-onliness.

Next, we present a minimal fairness condition. It states that whenever an agent has equal division as his preference's peak, the rule must assign it to him.

Equal division guarantee: For each $(R, \Omega) \in \mathcal{D}$ and each $i \in N$ such that $\frac{\Omega}{n} \in p(R_i)$, we have $\varphi_i(R) I_i \frac{\Omega}{n}$.

Another well-known fairness property states that agents with the same preferences should be assigned indifferent allocations.

Symmetry: For each $(R, \Omega) \in \mathcal{D}$ and each $\{i, j\} \subseteq N$ such that $R_i = R_j$, $\varphi_i(R, \Omega) I_i \varphi_j(R, \Omega)$.

²This property is weaker than the "peak-only" property, that has been imposed in a number of axiomatic studies. See [Section 6](#) for more details.

2.2 Obvious manipulations

The notion of obvious manipulation is introduced by [Trojan and Morrill \(2020\)](#). They try to single out those manipulations that are easily identifiable by the agents. A manipulation is obvious if the worst possible outcome under the manipulation is strictly better than the worst possible outcome under truth-telling.³ When the set of alternatives is infinite and preferences allow indifferences, worst possible outcomes may not be well-defined and a more general definition is necessary in this case. We say that a manipulation is obvious if each possible outcome under the manipulation is strictly better than some possible outcome under truth-telling.

Before we present the formal definition, we introduce some notation. Given a rule φ defined on $\mathcal{E}_{\mathcal{D}}$ and $(R_i, \Omega) \in \mathcal{D} \times \mathbb{R}_{++}$, the **option set** attainable with (R_i, Ω) at φ is

$$O^\varphi(R_i, \Omega) = \left\{ \varphi_i(R_i, R_{-i}) \in [0, \Omega] \text{ such that } R_{-i} \in \mathcal{D}^{n-1} \right\}.^4$$

Definition 1 *Let φ be a rule defined on $\mathcal{E}_{\mathcal{D}}$ and let $R_i, R'_i \in \mathcal{D}$ and $\Omega \in \mathbb{R}_{++}$ be such that R'_i is a manipulation of φ at (R_i, Ω) . Then, R'_i is an **obvious manipulation** if for each $x' \in O^\varphi(R'_i, \Omega)$ there is $x \in O^\varphi(R_i, \Omega)$ such that $x' P_i x$. The rule φ is **not obviously manipulable (NOM)** if it does not admit any obvious manipulation.*

The next lemma shows that our definition generalizes [Trojan and Morrill \(2020\)](#)'s definition. In fact, both definitions are equivalent whenever worst possible outcomes exist. Given R_i and $Y \subseteq [0, \Omega]$, denote by $W(R_i, Y)$ the element in Y (if any) such that $x P_i W(R_i, Y)$ for each $x \in Y \setminus \{W(R_i, Y)\}$.

Lemma 1 *Let φ be a rule defined on $\mathcal{E}_{\mathcal{D}}$ and let $R_i, R'_i \in \mathcal{D}$ and $\Omega \in \mathbb{R}_{++}$ be such that R'_i is a manipulation of φ at (R_i, Ω) . Assume that $W(R_i, O^\varphi(R_i, \Omega))$ and $W(R_i, O^\varphi(R'_i, \Omega))$ exist. Then, R'_i is an obvious manipulation if and only if $W(R_i, O^\varphi(R'_i, \Omega)) P_i W(R_i, O^\varphi(R_i, \Omega))$.*

³[Trojan and Morrill \(2020\)](#) also consider that a manipulation is obvious when the best possible outcome under the manipulation is strictly better than the best possible outcome under truth-telling. However, under EFF, the best possible outcome under truth-telling for an agent is always a peak alternative for that agent in our model. Therefore, no manipulation becomes obvious by considering best possible outcomes and so we omit this in our definition.

⁴[Barbera and Peleg \(1990\)](#) were the first to use option sets in the context of preference aggregation.

Proof. (\implies) Let R'_i be an obvious manipulation of φ at (R_i, Ω) . As $W(R_i, O^\varphi(R'_i, \Omega)) \in O^\varphi(R'_i, \Omega)$, there is $x \in O^\varphi(R_i, \Omega)$ such that $W(R_i, O^\varphi(R_i, \Omega))P_i x$. Then, by definition of $W(R_i, O^\varphi(R_i, \Omega))$,

$$W(R_i, O^\varphi(R'_i, \Omega))P_i W(R_i, O^\varphi(R_i, \Omega)) \quad (1)$$

(\impliedby) Assume that (1) holds. We will see that R'_i is an obvious manipulation of φ at (R_i, Ω) . Let $x' \in O^\varphi(R'_i, \Omega)$. Then, by definition of $W(R_i, O^\varphi(R'_i, \Omega))$ and (1),

$$x' R_i W(R_i, O^\varphi(R'_i, \Omega))P_i W(R_i, O^\varphi(R_i, \Omega)).$$

As $W(R_i, O^\varphi(R_i, \Omega)) \in O^\varphi(R_i, \Omega)$, the proof is complete. \square

3 Simple rules

In subsection 3.1, the formal definition of simple rule and our two main characterizations are presented. In subsection 3.2, the independence of the axioms involved in the characterizations is discussed.

3.1 Definition and characterizations

Agents i 's preference $R_i \in \mathcal{U}$ is **single-peaked** if $p(R_i)$ is a singleton and, for each pair $\{x_i, x'_i\} \subseteq \mathbb{R}_+$, we have $x_i P_i x'_i$ as long as either $x'_i < x_i \leq p(R_i)$ or $p(R_i) \leq x_i < x'_i$ holds. Denote by \mathcal{SP} the domain of all such preferences. We introduce a useful property that is equivalent to efficiency on $\mathcal{E}_{\mathcal{SP}}$.

Same-sidedness: For each $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$,

- (i) $\sum_{j \in N} p(R_j) \geq \Omega$ implies $\varphi_i(R) \leq p(R_i)$ for each $i \in N$, and
- (ii) $\sum_{j \in N} p(R_j) \leq \Omega$ implies $\varphi_i(R) \geq p(R_i)$ for each $i \in N$.

Remark 1 Let φ be a rule defined on $\mathcal{E}_{\mathcal{SP}}$. Then, φ is efficient if and only if it is same-sided.

Given $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$, let $z(R, \Omega) = \sum_{j \in N} p(R_j) - \Omega$. If $z(R, \Omega) \geq 0$ we say that economy (R, Ω) has **excess demand** whereas if $z(R, \Omega) < 0$ we say that economy (R, Ω) has

excess supply.⁵ Given economy $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$, agent $i \in N$ is **simple** if either $z(R, \Omega) \geq 0$ and $p(R_i) < \frac{\Omega}{n}$ or $z(R, \Omega) \leq 0$ and $p(R_i) > \frac{\Omega}{n}$. Let $N^*(R, \Omega)$ denote the set of simple agents of economy (R, Ω) .

Given $x, y, z \in \mathbb{R}$, we say that x is **between y and z** if $y \leq z$ and $x \in [y, z]$ or $z \leq y$ and $x \in [z, y]$.

In economies with excess demand, simple rules fully satiate agents whose peak amount is no greater than equal division and assign to each remaining agent an amount between equal division and his peak. Symmetrically, in economies with excess supply, simple rules fully satiate agents whose peak is no smaller than equal division and assign to each remaining agent an amount between his peak and equal division. Formally,

Definition 2 An own-peak-only rule φ defined on $\mathcal{E}_{\mathcal{SP}}$ is **simple** if for each $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$ and each $i \in N$,

(i) $\varphi_i(R, \Omega)$ is between $\frac{\Omega}{n}$ and $p(R_i)$ for each $i \in N$, and

(ii) $\varphi_i(R, \Omega) = p(R_i)$ if $i \in N^*(R, \Omega)$.

A prominent member of the family of simple rules is the uniform rule, first characterized by [Sprumont \(1991\)](#).

Uniform rule, u : For each $(R, \Omega) \in \mathcal{E}^N$ and each $i \in N$,

$$u_i(R, \Omega) = \begin{cases} \min\{p(R_i), \lambda\} & \text{if } z(R, \Omega) \geq 0 \\ \max\{p(R_i), \lambda\} & \text{if } z(R, \Omega) < 0 \end{cases}$$

where $\lambda \geq 0$ and solves $\sum_{j \in N} u_j(R, \Omega) = \Omega$.

The following result shows that simple rules cannot be obviously manipulated.

Theorem 1 A simple rule defined on $\mathcal{E}_{\mathcal{SP}}$ satisfies NOM.

Proof. Let φ be a simple rule defined on $\mathcal{E}_{\mathcal{SP}}$, and let $i \in N$ and $(R_i, \Omega) \in \mathcal{SP} \times \mathbb{R}_{++}$.

$$\text{Claim: } O^\varphi(R_i, \Omega) = \begin{cases} \left[\frac{\Omega}{n}, p(R_i) \right] & \text{if } p(R_i) > \frac{\Omega}{n} \\ \left[p(R_i), \frac{\Omega}{n} \right] & \text{if } p(R_i) \leq \frac{\Omega}{n} \end{cases}$$

⁵We deliberately include the **balanced** case where $z(R, \Omega) = 0$ in the excess demand case. No confusion will arise.

Assume that $\frac{\Omega}{n} < p(R_i)$ (the other case is symmetric). That $O^\varphi(R_i, \Omega) \subseteq [\frac{\Omega}{n}, p(R_i)]$ is clear by the definition of the rule. Now, let $x \in [\frac{\Omega}{n}, p(R_i)]$. We will prove that $x \in O^\varphi(R_i, \Omega)$. Consider $R_{-i} \in \mathcal{SP}^{n-1}$ such that $p(R_j) = \frac{\Omega}{n} - \frac{x - \frac{\Omega}{n}}{n-1}$ for each $j \in N \setminus \{i\}$. Then, as $\frac{\Omega}{n} \leq x$, $p(R_j) \leq \frac{\Omega}{n} \leq x$ for all $j \in N \setminus \{i\}$ and $\sum_{j \in N \setminus \{i\}} p(R_j) + x = \Omega$. By definition of the rule, $\varphi_j(R, \Omega) = p(R_j)$ for all $j \in N \setminus \{i\}$. Then, by feasibility, $\varphi_i(R, \Omega) = x$. Therefore, $[\frac{\Omega}{n}, p(R_i)] \subseteq O^\varphi(R_i, \Omega)$. This completes the proof of the Claim.

To see that φ satisfies NOM, assume that R'_i is a manipulation of φ at (R_i, Ω) . By the Claim, $\frac{\Omega}{n} \in O^\varphi(R'_i)$. Furthermore, by the Claim and single-peakedness, $x R_i \frac{\Omega}{n}$ for each $x \in O^\varphi(R_i)$. Then, R'_i is not an obvious manipulation and so φ is NOM. \square

Next, we present our first characterization result.

Theorem 2 *An own-peak-only rule defined on $\mathcal{E}_{\mathcal{SP}}$ satisfies efficiency, the equal division guarantee, and NOM if and only if it is a simple rule.*

Proof. (\Leftarrow) Let φ be a simple rule on $\mathcal{E}_{\mathcal{SP}}$. Then, φ is trivially own-peak-only and satisfies the equal division guarantee. By definition, φ satisfies same-sidedness. Therefore, by Remark 1, φ is efficient. Lastly, by Theorem 1, φ satisfies NOM.

(\Rightarrow) Let φ be an own-peak-only rule defined on $\mathcal{E}_{\mathcal{SP}}$ that satisfies efficiency, the equal division guarantee, and NOM. We now prove that φ is simple. Let $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$ and assume that $z(R, \Omega) \geq 0$. Let $i \in N$. By efficiency, $\varphi_i(R, \Omega) \leq p(R_i)$. There are two cases to consider:

1. $p(R_i) \leq \frac{\Omega}{n}$. We need to show that $\varphi_i(R, \Omega) \geq p(R_i)$ also holds. Assume, by way of contradiction, that $\varphi_i(R, \Omega) < p(R_i)$. By own-peak-onliness, we can assume that R_i is such that $\frac{\Omega}{n} P_i \varphi_i(R, \Omega)$. Let $R'_i \in \mathcal{SP}$ be such that $p(R'_i) = \frac{\Omega}{n}$. By the equal division guarantee, $\varphi_i(R'_i, R_{-i}, \Omega) = \frac{\Omega}{n}$. Hence,

$$\varphi_i(R'_i, R_{-i}, \Omega) = \frac{\Omega}{n} P_i \varphi_i(R, \Omega) \tag{2}$$

and R'_i is a manipulation of φ at (R_i, Ω) . Furthermore, by the equal division guarantee, $O^\varphi(R'_i, \Omega) = \left\{ \frac{\Omega}{n} \right\}$. As $\varphi_i(R, \Omega) \in O^\varphi(R_i, \Omega)$, by (2), R'_i is an obvious manipulation of φ at (R_i, Ω) , contradicting that φ is NOM. Thus, $\varphi_i(R, \Omega) \geq p(R_i)$ and, since by hypothesis $\varphi_i(R, \Omega) \leq p(R_i)$, we have $\varphi_i(R, \Omega) = p(R_i)$.

2. $p(R_i) > \frac{\Omega}{n}$. We need to show that $\varphi_i(R, \Omega) \geq \frac{\Omega}{n}$. Assume, by way of contradiction, that $\varphi_i(R, \Omega) < \frac{\Omega}{n}$. By single-peakedness and own-peak-onliness, $\frac{\Omega}{n} P_i \varphi_i(R, \Omega)$ also holds here, and the proof follows an argument similar to that of the previous case.

We have proved that if $p(R_i) \leq \frac{\Omega}{n}$, $\varphi_i(R, \Omega) = p(R_i)$; whereas if $p(R_i) > \frac{\Omega}{n}$, $\varphi_i(R, \Omega) \in [\frac{\Omega}{n}, p(R_i)]$. To complete the proof that φ is a simple rule, a symmetric argument can be used for the case $z(R, \Omega) < 0$. \square

Simple rules only meet a minimal fairness requirement: whenever an agent has equal division as his peak consumption, the rule must assign it to him. Despite this requirement, simple rules can be quite unfair and fail symmetry. If we want to avoid this unfairness adding symmetry to the picture, the following corollary follows.⁶

Corollary 1 *An own-peak-only rule defined on \mathcal{E}_{SP} satisfies efficiency, the equal division guarantee, symmetry, and NOM if and only if it is a symmetric and simple rule.*

Ching (1994) shows that if we require strategy-proofness instead of NOM in Corollary 1, the unique rule that prevails is the uniform rule. Therefore, Corollary 1 evidences how the family of characterizing rules enriches when NOM is sufficient to describe the strategic behavior of the agents.

An interesting member of this family is the simple rule that applies the constrained equal distance criterion to non-simple agents. If we use the distance between assignments and peak amounts as a measure of the sacrifice made by the agents at an allocation, this criterion tries to equate that sacrifice among agents as much as possible. It is easy to see that this criterion is incompatible with the definition of simple rule when all agents are considered. However, it is still relevant when applied only to non-simple agents.

CED simple rule, φ^{CED} : For each $(R, \Omega) \in \mathcal{E}^N$ and each $i \in N$,

(i) if $i \in N^*(R, \Omega)$, $\varphi_i^{CED}(R, \Omega) = p(R_i)$,

(ii) if $i \in N \setminus N^*(R, \Omega)$,

$$\varphi_i^{CED}(R, \Omega) = \begin{cases} \max\{\frac{\Omega}{n}, p(R_i) - d\} & \text{if } z(R, \Omega) \geq 0 \\ \min\{\frac{\Omega}{n}, p(R_i) + d\} & \text{if } z(R, \Omega) < 0 \end{cases}$$

where $d \geq 0$ and solves $\sum_{j \in N} \varphi_j^{CED}(R, \Omega) = \Omega$.

⁶As we will see in Subsection 3.2, the equal division guarantee and symmetry are independent axioms.

3.2 Independence of axioms

Next, we present some examples that prove the independence of the axioms in Theorem 2 and Corollary 1.

- Let $\tilde{\varphi} : \mathcal{E}_{\mathcal{SP}} \rightarrow \mathbb{R}_+^N$ be such that, for each $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$ and each $i \in N$, $\varphi_i(R, \Omega) = \frac{\Omega}{n}$. Then $\tilde{\varphi}$ satisfies all properties but efficiency.
- Let $\varphi^* : \mathcal{E}_{\mathcal{SP}} \rightarrow \mathbb{R}_+^N$ be such that, for each $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$ and each $i \in N$,

$$\varphi^*(R, \Omega) = \begin{cases} (p(R_1), p(R_2), 0, \dots, 0) & \text{if } p(R_1) + p(R_2) = \Omega \text{ and} \\ & p(R_j) \notin \{p(R_1), p(R_2)\} \text{ for each } j \in N \setminus \{1, 2\} \\ u(R, \Omega) & \text{otherwise} \end{cases}$$

Observe that this rule does meet the equal division guarantee. It is clear that φ^* is own-peak-only, efficient, and symmetric. Now, we prove that φ^* also satisfies NOM. Let $i \in N$ and $(R_i, \Omega) \in \mathcal{SP} \times \mathbb{R}_+$. Notice that $O^{\varphi^*}(R_i, \Omega) = O^u(R_i, \Omega) \cup \{0\}$ if $i \in N \setminus \{1, 2\}$ and $O^{\varphi^*}(R_i, \Omega) = O^u(R_i, \Omega)$ if $i \in \{1, 2\}$. Then, the fact that φ^* does not have obvious manipulations for agents in $N \setminus \{1, 2\}$ follows a similar argument as in Theorem 1 using the fact that u is a simple rule. Then, φ^* satisfies all properties but the equal division guarantee.

- Let $\bar{\varphi} : \mathcal{E}_{\mathcal{SP}} \rightarrow \mathbb{R}_+^N$ be such that, for each $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$,

$$\bar{\varphi}(R, \Omega) = \begin{cases} (\frac{1}{3}\Omega, \frac{2}{3}\Omega, 0, \dots, 0) & \text{if } p(R_1) = p(R_2) = \Omega \text{ and} \\ & p(R_j) = 0 \text{ for each } j \in N \setminus \{1, 2\} \\ u(R, \Omega) & \text{otherwise} \end{cases}$$

Observe that $\bar{\varphi}$ is simple. Thus, it satisfies own-peak-onliness, efficiency, the equal division guarantee, and NOM. However, it does not satisfy symmetry.

- Let $\hat{\varphi} : \mathcal{E}_{\mathcal{SP}} \rightarrow \mathbb{R}_+^N$ be such that, for each $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$,

(i) if $z(R, \Omega) \geq 0$, $\hat{\varphi}(R, \Omega) = u(R, \Omega)$.

(ii) if $z(R, \Omega) < 0$, let $\hat{N} = \{i \in N : p(R_i) = \min_{j \in N} \{p(R_j)\}\}$. Then, for each $i \in N$,

$$\hat{\varphi}_i(R, \Omega) = \begin{cases} p(R_i) & \text{if } i \in N \setminus \hat{N} \\ \lambda & \text{if } i \in \hat{N} \end{cases}$$

where $\lambda \geq 0$ and solves $\sum_{j \in N} \hat{\varphi}_j(R, \Omega) = \Omega$.

Observe that $\widehat{\varphi}$ is not simple because it could be that $p(R_i) < \frac{\Omega}{n} < \widehat{\varphi}_i(R)$ for some $i \in \widehat{N}$. It is easy to see that $\widehat{\varphi}$ is own-peak-only, efficient, and symmetric. If $z(R, \Omega) < 0$ and $p(R_i) = \frac{\Omega}{n}$, then $i \in N \setminus \widehat{N}$. Therefore, $\widehat{\varphi}$ satisfies the equal division guarantee. This implies, by Theorem 2, that $\widehat{\varphi}$ satisfies all properties but NOM.

- Let $\underline{\varphi} : \mathcal{E}_{\mathcal{SP}} \rightarrow \mathbb{R}_+^N$ be such that, for each $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$ and each $i \in N$,

$$\underline{\varphi}_i(R, \Omega) = \begin{cases} u_i(R_1^0, R_{-1}, \Omega) & \text{if } z(R_{-1}, \Omega) \geq 0, 0P_1 \frac{\Omega}{n}, \text{ and} \\ & p(R_1) < p(R_j) \text{ for each } j \in N \setminus \{1\} \\ u_i(R, \Omega) & \text{otherwise} \end{cases}$$

where $R_1^0 \in \mathcal{SP}$ is such that $p(R_1^0) = 0$.

Observe that $\underline{\varphi}$ is not simple because it is not own-peak-only and it could be that $\varphi_1(R, \Omega) < p(R_1) < \frac{\Omega}{n}$. It is clear that $\underline{\varphi}$ satisfies the equal division guarantee and symmetry. To see that $\underline{\varphi}$ is efficient, let $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$ be such that $z(R_{-1}, \Omega) \geq 0$, $0P_1 \frac{\Omega}{n}$, and $p(R_1) < p(R_j)$ for each $j \in N \setminus \{1\}$. Then, by the definition of u and the fact that u satisfies same-sideness, $\varphi_1(R, \Omega) = 0 \leq p(R_1)$ and $\varphi_i(R, \Omega) \leq p(R_i)$ for each $i \in N \setminus \{1\}$. Then, $\varphi_i(R, \Omega) \leq p(R_i)$ for each $i \in N$. Since in any other case $\underline{\varphi}$ operates as u and u satisfies same-sideness, $\underline{\varphi}$ also satisfies this property. Thus, by Remark 1, $\underline{\varphi}$ is efficient. Now, we prove that $\underline{\varphi}$ satisfies NOM. For each $i \in N \setminus \{1\}$ and each $(R_i, \Omega) \in \mathcal{SP} \times \mathbb{R}_+$, $O^{\underline{\varphi}}(R_i, \Omega) = O^u(R_i, \Omega)$. Then, the fact that $\underline{\varphi}$ does not have obvious manipulations for agents in $N \setminus \{1\}$ follows the same arguments used in the proof of Theorem 1 and the fact that u is a simple rule. Let us next consider agent 1. First, let $(R_1, \Omega) \in \mathcal{SP} \times \mathbb{R}_+$ be such that $\frac{\Omega}{n} R_1 0$. Then, $O^{\underline{\varphi}}(R_1, \Omega) = O^u(R_1, \Omega)$. Again, the fact that $\underline{\varphi}$ does not have obvious manipulations for agent 1 at (R_1, Ω) follows the same arguments used in the proof of Theorem 1 together with the facts that u is a simple rule and that $\frac{\Omega}{n} \in O^{\varphi}(R'_1, \Omega)$ for each $R'_1 \in \mathcal{SP}$. Second, let $(R_1, \Omega) \in \mathcal{SP} \times \mathbb{R}_+$ be such that $0P_1 \frac{\Omega}{n}$. Then, by single-peakedness of R_1 , $p(R_1) < \frac{\Omega}{n}$. By definition of $\underline{\varphi}$, $O^{\underline{\varphi}}(R_1, \Omega) = O^u(R_1^0, \Omega) = [0, \frac{\Omega}{n}]$. Using single-peakedness again and the fact that $0P_1 \frac{\Omega}{n}$, it follows that $xR_1 \frac{\Omega}{n}$ for each $x \in O^{\varphi}(R_1, \Omega)$. Furthermore, as $\frac{\Omega}{n} \in O^{\varphi}(R'_1, \Omega)$ for each $R'_1 \in \mathcal{SP}$, agent 1 does not have obvious manipulations of $\underline{\varphi}$ at (R_1, Ω) . Therefore, $\underline{\varphi}$ is NOM. We conclude that $\underline{\varphi}$ satisfies all properties but own-peak-onliness.

4 Further characterizations involving NOM

Insisting on the NOM property, in this section we provide new characterizations of rules in which NOM together with other relevant properties are invoked. These considerations lead to three interesting variants of Theorem 2. The first one involves the peaks-only property, the second one is an adaptation of Theorem 2 to economies with individual endowments, and the third one involves a monotonicity condition concerning the peaks of the preference profile.

4.1 Peaks-onliness

The following strong form of the *own-peak-only* property requires that if two profiles of preferences have the same peaks, then the allocations recommended by the rule in both profiles are the same.⁷

Peaks-only: For each $e = (R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$ and each $R' \in \mathcal{SP}^n$ such that $p(R'_i) = p(R_i)$ for each $i \in N$, $\varphi(R, \Omega) = \varphi(R', \Omega)$.

If we replace own-peak-onliness by peaks-onliness in the definition of simple rules in Section 3, then the following variant of Theorem 2 is obtained.

Theorem 3 *A rule defined on $\mathcal{E}_{\mathcal{SP}}$ satisfies peaks-onliness, efficiency, the equal division guarantee, and NOM if and only if it is a peaks-only and simple rule.*

The analysis of the independence of axioms involved is similar to that of Theorem 2 (see Subsection 3.2). For Theorem 3, rule $\underline{\varphi}$ is used to show a rule that satisfies all properties but peaks-onliness.

4.2 Endowments guarantee

In some situations, it is more appropriate to assume that instead of a social endowment $\Omega \in \mathbb{R}_+$, there is a profile $\omega = (\omega_i)_{i \in N} \in \mathbb{R}_+^N$ where, for each $i \in N$, ω_i denotes agent i 's individual endowment of the non-disposable commodity and $\sum_{j \in N} \omega_j = \Omega$. In such cases, it is natural to replace the equal division guarantee with the following property.

⁷Although similar to own-peak-onliness, peaks-onliness does *not* follow from efficiency and strategy-proofness.

Endowments guarantee: For each $(R, \omega) \in \mathcal{E}_{\mathcal{D}}$ and each $i \in N$ such that $\omega_i \in p(R_i)$, we have $\varphi_i(R)I_i\omega_i$.

If, for each $i \in N$, we replace $\frac{\Omega}{n}$ with ω_i in the definition of simple agent, we can define a **simple reallocation rule** as one that assigns: (i) to each agent, an amount between his endowment and his peak and, in particular, (ii) to each simple agent, his peak amount. Then, for economies with individual endowments, the following variant of Theorem 2 is obtained.

Theorem 4 *An own-peak-only rule defined on $\mathcal{E}_{\mathcal{SP}}$ satisfies efficiency, the endowments guarantee, and NOM if and only if it is a simple reallocation rule.*

The analysis of the independence of axioms involved is similar to that of Theorem 2 (see Subsection 3.2).

4.3 Peak monotonicity

The following monotonicity property says that if the peak of one agent is greater than or equal to the peak of another agent, then the first agent must be assigned an amount greater than or equal to that of the second agent.

Peak monotonicity: For each $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$ and each $\{i, j\} \subseteq N$ with $i \neq j$, $p(R_i) \leq p(R_j)$ implies $\varphi_i(R, \Omega) \leq \varphi_j(R, \Omega)$.

It is clear that peak monotonicity implies symmetry. Furthermore, we next prove that peak monotonicity together with NOM imply the equal division guarantee for an own-peak-only rule defined on $\mathcal{E}_{\mathcal{SP}}$.

Lemma 2 *Any own-peak-only rule defined on $\mathcal{E}_{\mathcal{SP}}$ that satisfies peak monotonicity and NOM meets the equal division guarantee.*

Proof. Let φ be an own-peak-only rule defined on $\mathcal{E}_{\mathcal{SP}}$ that satisfies peak monotonicity and NOM. Assume, on contradiction, that there are $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$ and $i \in N$ such that $p(R_i) = \frac{\Omega}{n}$ and $\varphi_i(R, \Omega) \neq \frac{\Omega}{n}$. Consider the case $\varphi_i(R, \Omega) < \frac{\Omega}{n}$ (the other one is symmetric). Let $R'_i \in \mathcal{SP}$ be such that $p(R'_i) = \infty$. By peak monotonicity, $\varphi_i(R'_i, R_{-i}, \Omega) \geq \varphi_j(R'_i, R_{-i}, \Omega)$ for each $j \in N \setminus \{i\}$. Then, by feasibility, $\varphi_i(R'_i, R_{-i}, \Omega) \geq \frac{\Omega}{n}$. By own-peak-onliness, we can assume that R_i is such that $\varphi_i(R'_i, R_{-i}, \Omega)P_i\varphi_i(R, \Omega)$. Then, R'_i is a manipulation of φ

at (R_i, Ω) . Furthermore, if $x \in O^\varphi(R'_i, \Omega)$, then $x = \varphi_i(R'_i, R'_{-i})$ for some $R'_{-i} \in \mathcal{SP}^{n-1}$. Then, by peak monotonicity and feasibility, $x \geq \frac{\Omega}{n}$. Therefore, by own-peak-onliness, we can assume that $x R_i \varphi_i(R)$. Therefore, R'_i is an obvious manipulation of φ at (R, Ω) , which is a contradiction. \square

Hence, under own-peak-onliness and NOM, peak monotonicity implies the fairness properties of symmetry and the equal division guarantee.⁸ By using Lemma 2 and a proof similar to that of Theorem 2, the following characterization is obtained.

Theorem 5 *An own-peak-only rule defined on $\mathcal{E}_{\mathcal{SP}}$ satisfies efficiency, peak monotonicity, and NOM if and only if it is a peak monotonic and simple rule.*

Let us check the independence of axioms involved in Theorem 5. Considering the rules defined in Subsection 3.2, the equal division rule $\tilde{\varphi}$ satisfies all properties but efficiency; by Lemma 2, rule φ^* satisfies all properties but peak monotonicity; and rule $\underline{\varphi}$ satisfies all properties but own-peak-onliness. Lastly, consider the following rule:

Constrained equal-distance rule, c : For each $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$ and each $i \in N$,

$$c_i(R, \Omega) = \begin{cases} \max\{p(R_i) - d, 0\} & \text{if } z(R, \Omega) \geq 0 \\ p(R_i) + d & \text{if } z(R, \Omega) < 0 \end{cases}$$

where $d \geq 0$ and solves $\sum_{j \in N} c_j(R, \Omega) = \Omega$.

This rule is own-peak-only by definition; it is same-sided, and therefore efficient; and it is also peak monotonic. However, since it may not assign the peak amount to some simple agent, it is not NOM.

Finally, Table 1 summarizes the different characterizations involving NOM presented in this study.

5 Maximal domain for NOM

Now, a relevant question that arises is whether the domain of preferences can be enlarged maintaining the compatibility of the properties.

⁸Remember that, as we saw in Section 3, symmetry and the equal division guarantee are independent properties.

	Th. 2	Cor. 1	Th. 3	Th. 4	Th. 5
Own-peak-onliness	+	+		+	+
Efficiency	+	+	+	+	+
Equal division guarantee	+	+	+		
NOM	+	+	+	+	+
Symmetry		+			
Peaks-onliness			+		
Endowments guarantee				+	
Peak monotonicity					+

Table 1: Characterizations of NOM rules.

Definition 3 A domain $\mathcal{D}^* \subseteq \mathcal{U}$ is **maximal for a list of properties** if (i) $\mathcal{SP} \subseteq \mathcal{D}^*$, (ii) there is a rule defined on $\mathcal{E}_{\mathcal{D}^*}$ satisfying the properties, and (iii) for each $\mathcal{E}_{\overline{\mathcal{D}}}$ with $\mathcal{D}^* \subsetneq \overline{\mathcal{D}} \subseteq \mathcal{U}$ there is no rule defined on $\mathcal{E}_{\overline{\mathcal{D}}}$ satisfying the same properties.

Given $i \in N$ and $R_i \in \mathcal{U}$, let $\underline{p}(R_i) = \inf p(R_i)$ and $\overline{p}(R_i) = \sup p(R_i)$. Agent i 's preference $R_i \in \mathcal{U}$ is **single-plateaued** if $p(R_i)$ is the interval $[\underline{p}(R_i), \overline{p}(R_i)]$ and, for each pair $\{x_i, x'_i\} \subseteq \mathbb{R}_+$, we have $x_i P_i x'_i$ as long as either $x'_i < x_i \leq \underline{p}(R_i)$ or $\overline{p}(R_i) \leq x_i < x'_i$ holds. Denote by $\mathcal{SP}\mathcal{L}$ the domain of all such preferences. Agent i 's preference $R_i \in \mathcal{U}$ is **convex** if $p(R_i)$ is the interval $[\underline{p}(R_i), \overline{p}(R_i)]$ and, for each pair $\{x_i, x'_i\} \subseteq \mathbb{R}_+$, we have $x_i R_i x'_i$ as long as either $x'_i < x_i \leq \underline{p}(R_i)$ or $\overline{p}(R_i) \leq x_i < x'_i$ holds. Denote by \mathcal{C} the domain of all such preferences.

Lemma 3 Let \mathcal{D} be such that $\mathcal{SP} \subseteq \mathcal{D} \subseteq \mathcal{U}$ and let φ be an own-peak-only rule defined on $\mathcal{E}_{\mathcal{D}}$ that satisfies the equal division guarantee and NOM. Let $i \in N$ and $(R, \Omega) \in \mathcal{E}_{\mathcal{D}}$. If $p(R_i)$ is a singleton, then $\varphi_i(R, \Omega)$ is between $p(R_i)$ and $\frac{\Omega}{n}$.

Proof. Let $\varphi : \mathcal{E}_{\mathcal{D}} \rightarrow \mathbb{R}_+$ be an own-peak-only rule that satisfies the equal division guarantee and NOM. Let $(R, \Omega) \in \mathcal{E}_{\mathcal{D}}$ and let $i \in N$ be such that $p(R_i) \leq \frac{\Omega}{n}$. We will prove that $\varphi_i(R, \Omega) \in [p(R_i), \frac{\Omega}{n}]$. By own-peak-onliness and the fact that $\mathcal{SP} \subseteq \mathcal{D}$, the proof that $\varphi_i(R, \Omega) \geq p(R_i)$ follows the same lines as Case 1 in part (\implies) of the proof of Theorem 2. Next, we need to show that $\varphi_i(R, \Omega) \leq \frac{\Omega}{n}$. Assume, by way of contradiction, that $\frac{\Omega}{n} < \varphi_i(R, \Omega)$. By own-peak-onliness and the fact that $\mathcal{SP} \subseteq \mathcal{D}$, we can assume w.l.o.g. that $\frac{\Omega}{n} P_i \varphi_i(R, \Omega)$. Let $R'_i \in \mathcal{SP}$ be such that $p(R'_i) = \frac{\Omega}{n}$. By the equal division

guarantee, R'_i is an obvious manipulation of φ at (R_i, Ω) . This contradicts the fact that φ satisfies NOM. Therefore, $\varphi_i(R, \Omega) \in [p(R_i), \frac{\Omega}{n}]$.

The case $\frac{\Omega}{n} > p(R_i)$ is symmetric, so we omit it. \square

Lemma 4 *Let \mathcal{D} be such that $\mathcal{SP} \subseteq \mathcal{D} \subseteq \mathcal{U}$. If there is an own-peak-only rule defined on $\mathcal{E}_{\mathcal{D}}$ that satisfies efficiency, the equal division guarantee, and NOM, then $\mathcal{D} \subseteq \mathcal{C}$.*

Proof. Let \mathcal{D} be such that $\mathcal{SP} \subseteq \mathcal{D} \subseteq \mathcal{U}$ and let φ be an own-peak-only rule defined on $\mathcal{E}_{\mathcal{D}}$ that is efficient, NOM, and meets the equal division guarantee. Assume, by way of contradiction, that there is $R_0 \in \mathcal{D} \setminus \mathcal{C}$. Then, there are $x, y, z \in \mathbb{R}_+$ such that $x < y < z$, xP_0y and zP_0y . Without loss of generality, we can assume that zR_0x (the case xR_0z is symmetric).

Let $x' = \max\{w \in [x, y] \text{ such that } wI_0x\}$ and $y' = \min\{w \in [y, z] \text{ such that } wI_0x\}$. As R_0 is continuous, both x' and y' are well defined, $x' < y < y'$, and $y'P_0w$ for each $w \in (x', y')$. Let $(R, \Omega) \in \mathcal{E}_{\mathcal{D}}$ be such that $\Omega = ny'$, $R_1 = R_0$, and $R_{-1} \in \mathcal{SP}^{n-1}$ is such that $p(R_j) \in (y', \frac{\Omega - x'}{n-1})$ and $\Omega R_j y'$ for each $j \in N \setminus \{1\}$. By Lemma 3, $\varphi_j(R, \Omega) \in [y', p(R_j)]$ for each $j \in N \setminus \{1\}$. Then, by feasibility and the fact that $p(R_j) < \frac{\Omega - x'}{n-1}$, it follows that $\varphi_1(R, \Omega) \in (x', y']$. If $\varphi_1(R, \Omega) \in (x', y')$, by the definitions of x' and y' , $y'P_1\varphi_1(R, \Omega)$. If we consider $R'_1 \in \mathcal{SP}$ be such that $p(R'_1) = \frac{\Omega}{n} = y'$, by the equal division guarantee R'_1 is an obvious manipulation of φ at (R_1, Ω) , contradicting that φ satisfies NOM. Therefore, $\varphi_1(R, \Omega) = y'$. Hence, again by Lemma 3 and feasibility, $\varphi_j(R, \Omega) = y'$ for each $j \in N$. Now, consider the feasible allocation in which agent 1 gets x' and the remaining agents get $\frac{\Omega - x'}{n-1}$. Since $x'I_1y'$ and $\frac{\Omega - x'}{n-1}P_jy'$ for each $j \in N \setminus \{1\}$, this new allocation Pareto improves upon $\varphi(R, \Omega)$, contradicting efficiency. We conclude that $\mathcal{D} \subseteq \mathcal{C}$. \square

Theorem 6 *Domain $\mathcal{SP}\mathcal{L}$ is a maximal domain for own-peak-onliness, efficiency, the equal division guarantee, and NOM.*

Proof. Let \mathcal{D} be a domain for the properties listed in the theorem such that $\mathcal{SP}\mathcal{L} \subseteq \mathcal{D}$. By Lemma 4, $\mathcal{D} \subseteq \mathcal{C}$. Assume, by way of contradiction, that there is $R_0 \in \mathcal{D} \setminus \mathcal{SP}\mathcal{L}$. Without loss of generality, we can assume that there are $x, y \in \mathbb{R}_+$ such that $x < y < \underline{p}(R_0)$ and xI_0y .

Assume further that there is a rule φ defined on $\mathcal{E}_{\mathcal{D}}$ that satisfies the properties listed in the Theorem. Consider $(R, \Omega) \in \mathcal{E}_{\mathcal{D}}$ with $\Omega = ny$, $R_1 = R_0$, and $R_{-1} \in \mathcal{SP}^{n-1}$ such

that $p(R_j) = y + \frac{y-x}{n-1}$ for each $j \in N \setminus \{1\}$. By Lemma 3, $\varphi_j(R, \Omega) \in [y, p(R_j)]$ for each $j \in N \setminus \{1\}$. Then, by feasibility, $\varphi_1(R, \Omega) \leq y$. If $\varphi_1(R, \Omega) < y$ we can assume, by own-peak-onliness and the fact that $\mathcal{SP}\mathcal{L} \subseteq \mathcal{D}$, that $yP_1\varphi_1(R, \Omega)$. Next, let us consider $R'_1 \in \mathcal{SP}$ such that $p(R'_1) = \frac{\Omega}{n} = y$. By the equal division guarantee, R'_1 is an obvious manipulation of φ at R_i , contradicting the fact that φ satisfies NOM. Therefore, $\varphi_1(R, \Omega) = y$. Hence, again by Lemma 3 and feasibility, $\varphi_j(R, \Omega) = y$ for each $j \in N$. Now, consider the feasible allocation in which agent 1 gets x and the remaining agents get $y + \frac{y-x}{n-1}$. Since xI_1y and $y + \frac{y-x}{n-1}P_jy$ for each $j \in N \setminus \{1\}$, this new allocation Pareto improves upon $\varphi(R, \Omega)$, contradicting efficiency. We conclude that $\mathcal{D} = \mathcal{SP}\mathcal{L}$.

The proof is completed by considering the extension of the uniform rule to domain $\mathcal{SP}\mathcal{L}$ defined in the proof of the Theorem in [Ching and Serizawa \(1998\)](#). As they observe, this rule satisfies efficiency and strategy-proofness on $\mathcal{SP}\mathcal{L}$. As NOM is a weakening of strategy-proofness, this rule satisfies NOM as well. Furthermore, it is easy to see that this rule is own-peak-only and meets the equal division guarantee.⁹ \square

6 Final Remarks

Some final remarks are in order. In this paper, we introduced the equal division guarantee as a minimal fairness requirement. We showed that it is not implied by symmetry. However, two well-known fairness properties, envy-freeness and the equal division lower bound, imply it for own-peak-only rules. Remember that a rule is envy-free if no agent prefers some other agent's allotment to his own and that a rule meets the equal division lower bound if, for each agent, the allotment recommended by the rule is at least as good

⁹We can also consider extensions of simple rules defined on $\mathcal{E}_{\mathcal{SP}}$ to domain $\mathcal{E}_{\mathcal{SP}\mathcal{L}}$. Before defining them, we need some notation. Given $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}\mathcal{L}}$, let $\underline{z}(R, \Omega) = \sum_{j \in N} \underline{p}(R_j) - \Omega$ and $\bar{z}(R, \Omega) = \sum_{j \in N} \bar{p}(R_j) - \Omega$. For $i \in N$ and $R_i \in \mathcal{SP}\mathcal{L}$, let $\underline{R}_i \in \mathcal{SP}$ be such that $p(\underline{R}_i) = \underline{p}(R_i)$ and let $\bar{R}_i \in \mathcal{SP}$ be such that $p(\bar{R}_i) = \bar{p}(R_i)$. An own-peak-only rule φ^* defined on $\mathcal{E}_{\mathcal{SP}\mathcal{L}}$ is an **extension** of simple rule φ defined on $\mathcal{E}_{\mathcal{SP}}$ if, for each $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}\mathcal{L}}$,

- (i) $\varphi_i^*(R, \Omega) = \varphi_i(\underline{R}, \Omega)$ if $\underline{z}(R, \Omega) \geq 0$,
- (ii) $\varphi_i^*(R, \Omega) = \varphi_i(\bar{R}, \Omega)$ if $\bar{z}(R, \Omega) \leq 0$,
- (iii) $\varphi_i^*(R, \Omega) \in [\underline{p}(R_i), \bar{p}(R_i)]$ if $\underline{z}(R, \Omega) < 0$ and $\bar{z}(R, \Omega) > 0$.

With similar arguments to those used in the proof of Theorem 2, it follows that these rules also are efficient, meet the equal division guarantee, and are NOM.

as equal division. Formally,

Envy-free: For each $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$ and each $\{i, j\} \subseteq N$ with $i \neq j$, $\varphi_i(R, \Omega)R_i \varphi_j(R, \Omega)$.

Equal division lower bound:¹⁰ For each $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$ and each $i \in N$, $\varphi_i(R, \Omega)R_i \frac{\Omega}{n}$.

Lemma 5 *For any own-peak-only rule defined on $\mathcal{E}_{\mathcal{SP}}$, envy-freeness or the equal division lower bound imply the equal division guarantee.*

Proof. Let φ be an own-peak-only rule defined on $\mathcal{E}_{\mathcal{SP}}$. First, assume that φ meets the equal division lower bound. Let $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$ and $i \in N$ be such that $p(R_i) = \frac{\Omega}{n}$. Then, $\varphi_i(R, \Omega)R_i \frac{\Omega}{n} = p(R_i)$ and thus $\varphi_i(R, \Omega) = p(R_i)$, implying that φ meets the equal division guarantee. Second, assume that φ is envy-free but does not meet the equal division guarantee. Then, there are $(R, \Omega) \in \mathcal{E}_{\mathcal{SP}}$ and $i \in N$ with $p(R_i) = \frac{\Omega}{n}$ such that $\varphi_i(R) \neq \frac{\Omega}{n}$. Consider the case $\varphi_i(R) > \frac{\Omega}{n}$ (the other one is symmetric). By feasibility, there is $j \in N \setminus \{i\}$ such that $\varphi_j(R) < \frac{\Omega}{n}$. As $p(R_i) = \frac{\Omega}{n}$, by own-peak-onliness, we can assume that R_i is such that $\varphi_j(R)P_i \varphi_i(R)$. Then, φ does not satisfy envy-freeness. This contradiction shows that φ meets the equal division guarantee. \square

Finally, concerning the maximality result of Section 5, it is worth mentioning that the single-plateaued domain remains maximal for the properties involved in the characterization of Corollary 1, i.e., when symmetry is also considered.

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¹⁰This property is also known as *individual rationality from equal division*.

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